# Déjà Q: Using Dual Systems to Revisit q-Type Assumptions 

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## Pairing-based cryptography: a brief history

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© Later assumptions: Subgroup Hiding [BGN05], Decision Linear, SXDH

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Even later assumptions: q-SDH, q-ADHSDH, q-EDBDH, q-SDH-III, q-SFP, "source group q-parallel BDHE," etc.

## Why are q-type assumptions worrisome?

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|  | IBE universe |  |  |
| :---: | :---: | :---: | :---: |
| Alice | Fred | Kate | Phil |
| Bob | George | Louise | Quentin |
| Charles | Hannah | Melissa | Rachel |
| Dora | Isabelle | Nicholas | Sarah |
| Ernie | Julian | Otis | Tristan |

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Extension to Dodis-Yampolskiy PRF [DY05]
*currently only in composite-order groups ${ }_{4}$

## Outline

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> q-Type assumptions


## Outline



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## Conclusions

## Properties of (bilinear) groups

Standard bilinear group: ( $\mathrm{N}, \mathrm{G}, \mathrm{H}, \mathrm{GT}, \mathrm{e}, \mathrm{g}, \mathrm{h}$ )

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$$
\begin{aligned}
& \text { Standard bilinear group: (N, G, H, GT, e, g, h) } \\
& \text { Group order; } \\
& \text { prime or composite } \\
& |G|=|H|=k N ;|G T|=\lambda N \\
& \mathrm{e}: \mathrm{G} \times \mathrm{H} \rightarrow \mathrm{G} \\
& \text { bilinearity: } e\left(g^{a}, h^{b}\right)=e(g, h)^{\text {ab }} \forall a, b \in Z / N Z \\
& \text { non-degeneracy: } e(x, y)=1 \quad \forall y \in H \Rightarrow x=1
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$$

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## Subgroup hiding

Composite-order bilinear group: $(\mathrm{N}, \mathrm{G}, \mathrm{G} \mathrm{t}, \mathrm{e}, \mathrm{g})$ where $\mathrm{N}=\mathrm{pq}$

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$\approx$


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\mathrm{g}_{1}{ }_{\mathrm{f}\left(\mathrm{x} 1, \ldots, \mathrm{x}_{\mathrm{c}}\right)}^{\bar{\equiv}}{ }_{\mathrm{g}_{2}{ }^{\mathrm{f}\left(\mathrm{x} 1, \ldots, \mathrm{x}_{\mathrm{c}}\right)}}
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Parameter hiding: elements correlated across subgroups are distributed identically to uncorrelated elements
$\bigcirc$ is independent from $\square$
$x_{x_{i}}$ mod p reveals nothing about $x_{i}$ mod $q$ (CRT)

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## SF keys don't decrypt SF ciphertexts!

ID queries
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Dual systems in three easy steps

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- $\mathrm{f}\left(\mathrm{x}_{1}, \ldots, \mathrm{x}_{\mathrm{c}}\right)$ : A needs to compute $\mathrm{e}(\mathrm{g}, \mathrm{h})^{\mathrm{f}\left(\mathrm{x}_{1}, \ldots, \mathrm{x}_{\mathrm{C}}\right)}$ (or distinguish it from random)


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- $\mathrm{f}\left(\mathrm{x}_{1}, \ldots, \mathrm{x}_{\mathrm{c}}\right)$ : A needs to compute $\mathrm{e}(\mathrm{g}, \mathrm{h})^{f\left(\mathrm{x}_{1}, \ldots, \mathrm{x}_{\mathrm{C}}\right)}$ (or distinguish it from random) uber(c,R,S,T,f) assumption: given (R,S,T) values, hard to compute/distinguish f


## Example uber-assumption: exponent q-SDH

exponent $\mathrm{q}-$ SDH [ZS-NSO4]: given $\left(\mathrm{g}, \mathrm{g}^{\mathrm{x}}, \ldots, \mathrm{g}^{\mathrm{x}}\right)$, distinguish $\mathrm{g}^{\mathrm{x}}{ }^{\mathrm{a}+1}$ from random

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- $\mathrm{T}=\langle 1\rangle$
- $f\left(x_{1}, \ldots, x_{c}\right): f(x)=x^{q+1}$
exponent q-SDH is $u b e r\left(1,<1,\left\{x^{i}\right\}>,<1>,<1>, x^{q+1}\right)$

Applying dual systems to exponent q-SDH
uber(c, $<1,\left\{x^{i}\right\}>,<1>,<1>$, $\left.^{q+1}\right)$

1. start with base scheme
2. transition to SF version
3. argue information is hidden ${ }_{14}$

## Applying dual systems to exponent q-SDH



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$g_{1} \sum r_{k x}, \ldots, g_{1} \sum r k x_{k} 9$

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## Applying dual systems to exponent q-SDH

```
uber(c,R,<1,{\mp@subsup{x}{}{i}}>,<1>,\mp@subsup{x}{}{q+1})->\operatorname{uber}(lc,<1,{\sum\mp@subsup{r}{k}{\prime}\mp@subsup{x}{k}{i}}>,<1>,<1>,\sum\mp@subsup{r}{k}{}\mp@subsup{x}{k}{}\mp@subsup{}{}{q+1})
```

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```
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```

$$
\left[\begin{array}{llll}
r_{1} & r_{2} & \ldots & r_{l}
\end{array}\right]\left[\begin{array}{ccccc}
1 & x & \cdot & x^{q} & x^{q+1} \\
1 & x_{2} & \cdot & x_{2}{ }^{q} & x_{2}{ }^{q+1} \\
\cdot & & \cdot & & \cdot \\
\cdot & & & \cdot & \cdot \\
1 & x_{\ell} & \cdot & x_{\ell}{ }^{q} & x_{l}{ }^{q+1}
\end{array}\right]
$$

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```
uber(c,R,<1,{\mp@subsup{x}{}{\prime}}>,<1>,\mp@subsup{x}{}{q+1})->\mathrm{ uber({c,<1,{ {rkxk}\mp@subsup{r}{k}{\prime}}>,<1>,<1>, \sum\mp@subsup{r}{k}{\prime}\mp@subsup{x}{k}{}\mp@subsup{}{}{q+1})
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\cdot & & \cdot & & \cdot \\
\cdot & & & \cdot & \cdot \\
1 & x_{\ell} & \cdot & x_{\ell}{ }^{q} & x_{\ell}{ }^{q+1}
\end{array}\right]=\left[\begin{array}{c}
y_{1} \\
y_{2} \\
\cdot \\
\cdot \\
y_{\ell}
\end{array}\right]
$$

So $A$ is really given

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r_{1} & r_{2} & \ldots & r_{l}
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1 & x & \cdot & x^{q} & x^{a+1} \\
1 & x_{2} & \cdot & x_{2}^{q} & x_{2}{ }^{q+1} \\
\cdot & & \cdot & & \cdot \\
\cdot & & & \cdot & \cdot \\
1 & x_{l} & \cdot & x_{l}{ }^{q} & x_{l}{ }_{l}+1
\end{array}\right]=\left[\begin{array}{c}
y_{1} \\
y_{2} \\
\cdot \\
\cdot \\
y_{l}
\end{array}\right]
$$

1. start with base scheme
2. transition to SF version
3. argue information is hidden

## Applying dual systems to exponent q-SDH

$$
\begin{array}{cccc}
{\left[\begin{array}{lllll}
r_{1} & r_{2} & \ldots & r_{l}
\end{array}\right]\left[\begin{array}{cccc}
1 & x & \cdot & x^{q} \\
x^{q+1} \\
1 & x_{2} & \cdot & x_{2}{ }^{q} \\
x_{2}{ }^{q+1} \\
\cdot & & \cdot & \\
\cdot & & \cdot \\
1 & x_{\ell} & \cdot & x_{\ell}{ }^{q} \\
x_{\ell}{ }_{l}^{q+1}
\end{array}\right]}
\end{array}=\left[\begin{array}{c}
y_{1} \\
y_{2} \\
\cdot \\
\cdot \\
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Consider set $S$ of $\ell$-sized sets; then $\mathbf{r}, \mathbf{y} \in \mathrm{S}$
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This is distributed uniformly random as well!
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## Applying dual systems to the uber-assumption

More generally, this is true if

| 1 | $\rho_{1}\left(\mathrm{X}_{11}, \ldots, \mathrm{x}_{10}\right)$ | $\rho_{q}\left(X_{11}, \ldots, X_{1 c}\right)$ | $f\left(x_{11}, \ldots, x_{10}\right)$ |
| :---: | :---: | :---: | :---: |
| 1 | $\rho_{1}\left(\mathrm{X}_{21}, \ldots, \mathrm{X}_{2 \mathrm{c}}\right)$ | $\rho_{q}\left(\mathrm{X}_{21}, \ldots, \mathrm{X}_{2 \mathrm{c}}\right)$ | $f\left(x_{21}, \ldots, x_{2 c}\right)$ |
| . | . |  |  |
| 1 | $\rho_{1}\left(\mathrm{X}_{\ell 1}, \ldots, \mathrm{X}_{\ell c}\right)$ | $\rho_{\mathrm{q}}\left(\mathrm{X}_{\ell 1}, \ldots, \mathrm{X}_{\ell c}\right)$ | $f\left(X_{\ell 1}, \ldots, X_{l c}\right)$ |

has linearly independent columns (or rows)

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2. transition to SF version
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## Outline



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...even when given a generator for $\bigcirc$

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$[q-S D H]\left(g, g^{x}, \ldots, g^{\times 9}, h^{x}\right) \rightarrow$ compute $\left(c, g^{1 / x+c}\right)$

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## limitation <br> Computational uber(c,R,S,T,f) holds if:

1. subgroup hiding and parameter hiding hold 2. $f$ is not a linear combination of $p_{i}$

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$$
\Longrightarrow \mathrm{sh} \xlongequal{\Longrightarrow} \equiv \mathrm{ph} \xlongequal{\Longrightarrow}
$$

This implies (for example) that q-SDH [BBO4] follows from subgroup hiding....
...and so does everything based on q-SDH (like Boneh-Boyen signatures)*
*when instantiated in asymmetric composite-order groups [BRS11]

## Reexamining the Dodis-Yampolskiy PRF

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\mathrm{f}(\mathrm{x})=\mathrm{u}^{1 / \mathrm{sk}+\mathrm{x}} \text { for fixed sk } \leftarrow R ; \mathrm{x} \in \mathrm{a}(\lambda)
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Theorem: Adv ${ }^{\text {prf }} \leq q \cdot A d v^{s g h}$

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\mathrm{f}(\mathrm{x})=\mathrm{u}^{-1 / \mathrm{sk}+\mathrm{x}} \text { for fixed sk } \leftarrow \mathcal{R} ; \mathrm{x} \in \mathrm{a}(\lambda)
$$

Theorem [DY05]: Adv ${ }^{v r f} \leq a(\lambda) \cdot A d v(\lambda)-D B D H$

## Theorem: Advorf $\leq q$ Advsgh

© pseudorandom function

-     - static assumption
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## Outline

> q-Type assumptions


## Conclusions

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Full version!: cs.ucsd.edu/~smeiklejohn/files/eurocrypt14a.pdf

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## Thanks! Any questions?

